

The SeqBin Constraint Revisited

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Australian Government

**Department of Communications,
Information Technology and the Arts**

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Outline

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1. The SeqBin constraint
2. Related work
3. New DC filtering algorithm

SeqBin

The SeqBin constraint
[IJCAI'11, Petit et al.]

SeqBin

Motivation

SeqBin

SeqBin is a meta-constraint useful for
(over-constrained) scheduling and rostering problems

SeqBin

SeqBin is a meta-constraint useful for (over-constrained) scheduling and rostering problems

- Change
- Smooth
- IncreasingNValue

SeqBin

$X_1, X_2, X_3, X_4, X_5, X_6, X_7$

SeqBin

$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$

$X_1, X_2, X_3, X_4, X_5, X_6, X_7$

SeqBin

a hard constraint

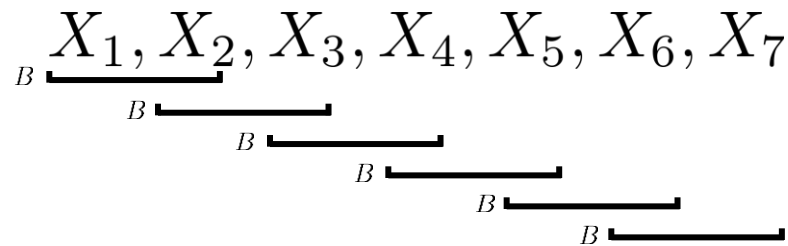
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SeqBin

a hard constraint


$$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$$

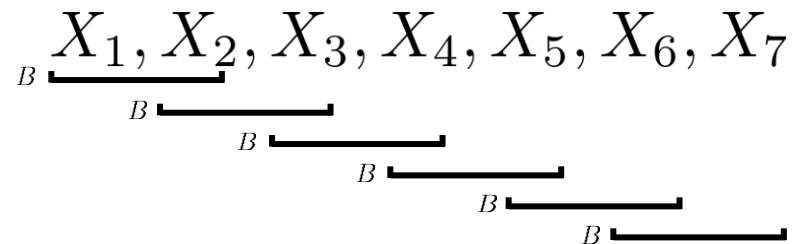


SeqBin

a soft constraint

a hard constraint


$$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$$


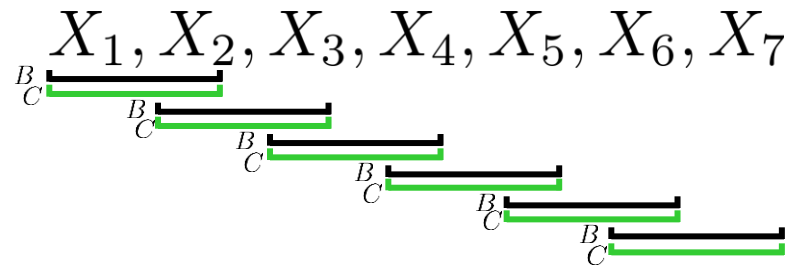


SeqBin

a soft constraint

a hard constraint

$$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$$




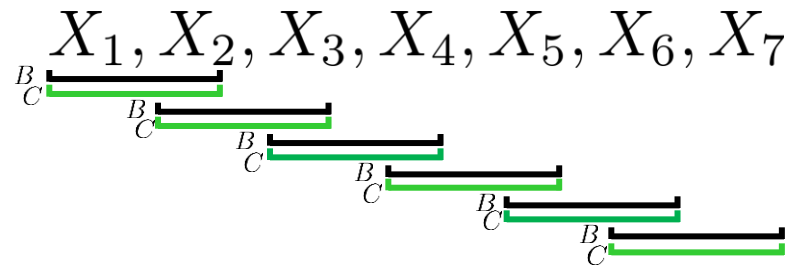
SeqBin

the nb of violated C

a soft constraint

a hard constraint

$$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$$



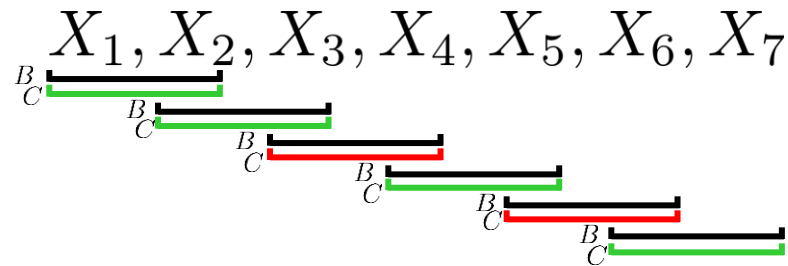
SeqBin

the nb of violated C

a soft constraint

a hard constraint

$$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$$



SeqBin

$\text{SeqBin}(N = 3, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$

SeqBin

$\text{SeqBin}(N = 3, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$

Variables

X_1

X_2

X_3

X_4

X_5

X_6

X_7

SeqBin

$\text{SeqBin}(N = 3, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$

Variables

X_1

X_2

X_3

X_4

X_5

X_6

X_7

Domains

1

1
0

1

1
0

1

1
0

1

SeqBin

$\text{SeqBin}(N = 3, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$

Variables

X_1

X_2

X_3

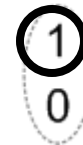
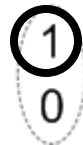
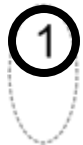
X_4

X_5

X_6

X_7

Domains



SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \checkmark \equiv \text{True}$)

Variables

X_1

X_2

X_3

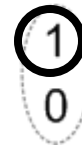
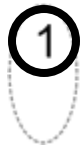
X_4

X_5

X_6

X_7

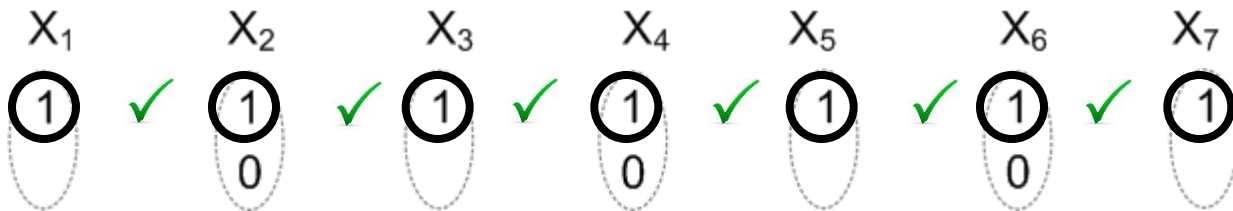
Domains



SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \stackrel{\checkmark}{=} \text{True}$)

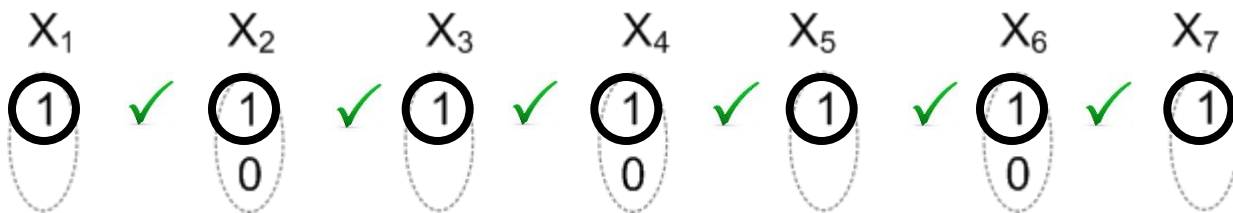
Variables



SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \stackrel{\checkmark}{=} \text{True}$)

Variables



$c(X) = 0$



SeqBin

$\text{SeqBin}(N = 3, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$

Variables

X_1

X_2

X_3

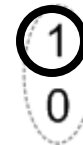
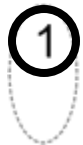
X_4

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Domains



SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \checkmark \equiv \text{True}$)

Variables

X_1

X_2

X_3

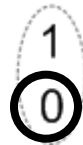
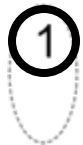
X_4

X_5

X_6

X_7

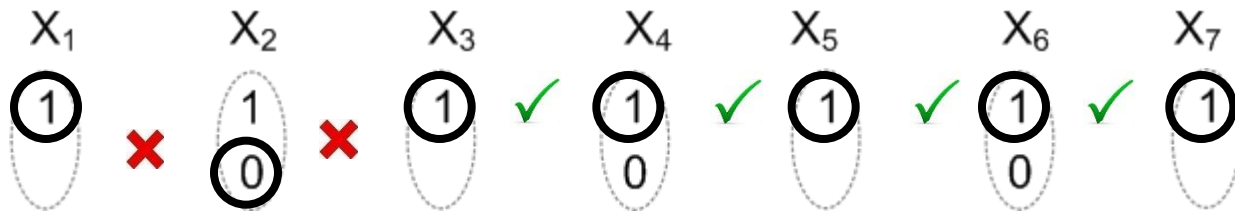
Domains



SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \stackrel{\checkmark}{=} \text{True}$)

Variables

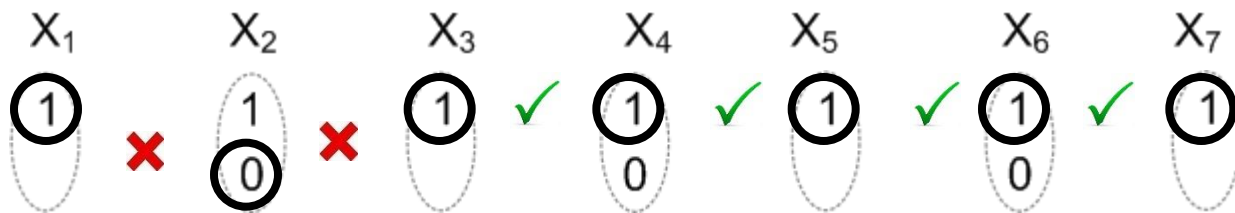


Domains

SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \checkmark \equiv \text{True}$)

Variables



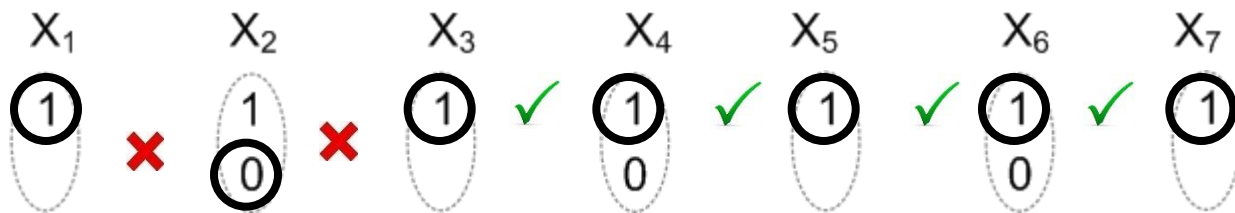
Domains

$$c(X) = 2$$

SeqBin

SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B \checkmark \equiv \text{True}$)

Variables



Domains

$$c(X) + 1 = 2 + 1 = N$$

SeqBin

Disentanglement detection

SeqBin

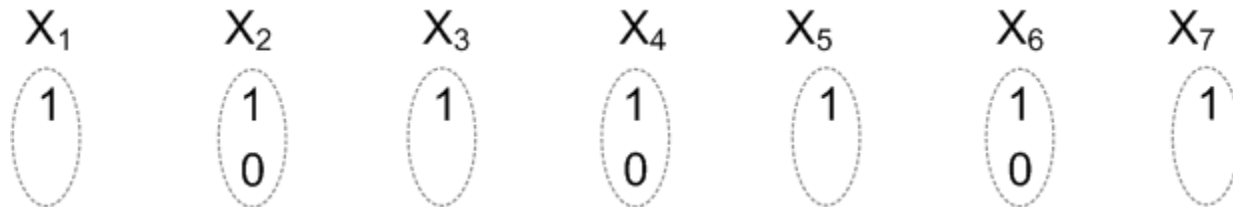
Is there a solution of SeqBin?

SeqBin (graph representation)

The SeqBin constraint
(equivalent representation)

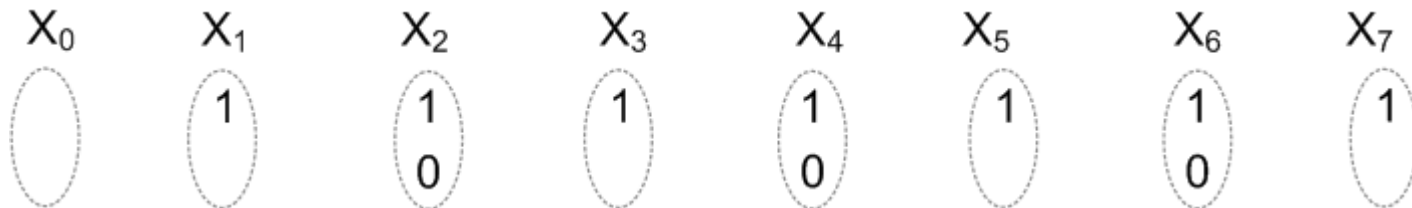
SeqBin (graph representation)

$\text{SeqBin}(N, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$



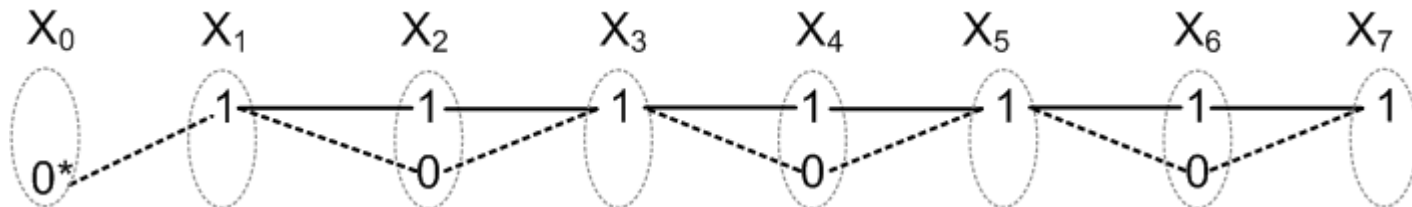
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SeqBin (graph representation)

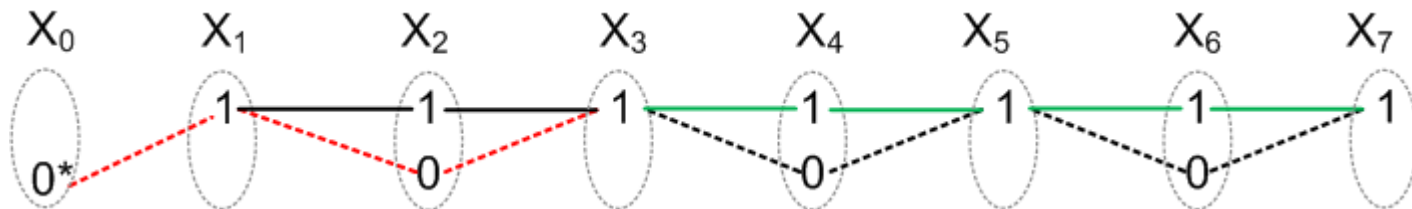
$\text{SeqBin}(N, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$



----- weight 1
——— weight 0

SeqBin (graph representation)

$\text{SeqBin}(N, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$



----- weight 1
——— weight 0

cost of the path is 3

SeqBin (graph representation)

There exists a bijection between assignments X of cost s that satisfy B and paths in the graph $G(V, E)$ of cost $s + 1$.

SeqBin

Is there a path in G with cost equal to N ?

SeqBin (DP-based propagator)

Straightforward approach

Dynamic programming based DC algorithm in $O(n^2d^2)$

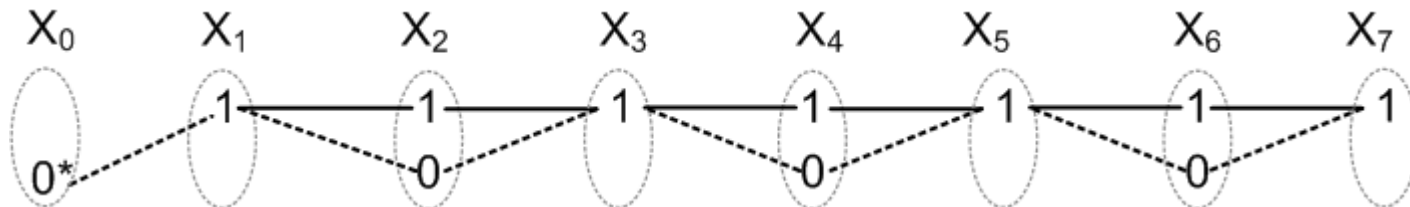
SeqBin (DP-based propagator)

For each vertex in G we compute costs of all possible paths from the sink node.

SeqBin (DP-based propagator)

----- weight 1
—— weight 0

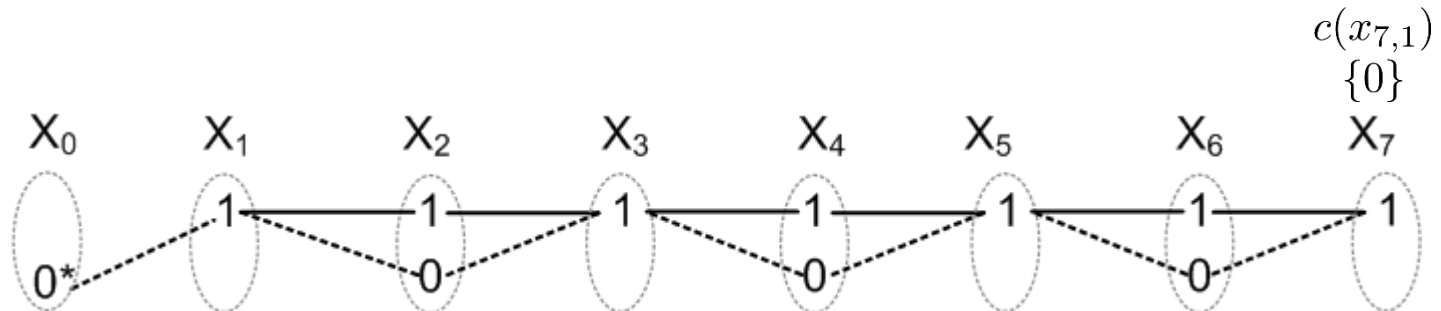
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SeqBin (DP-based propagator)

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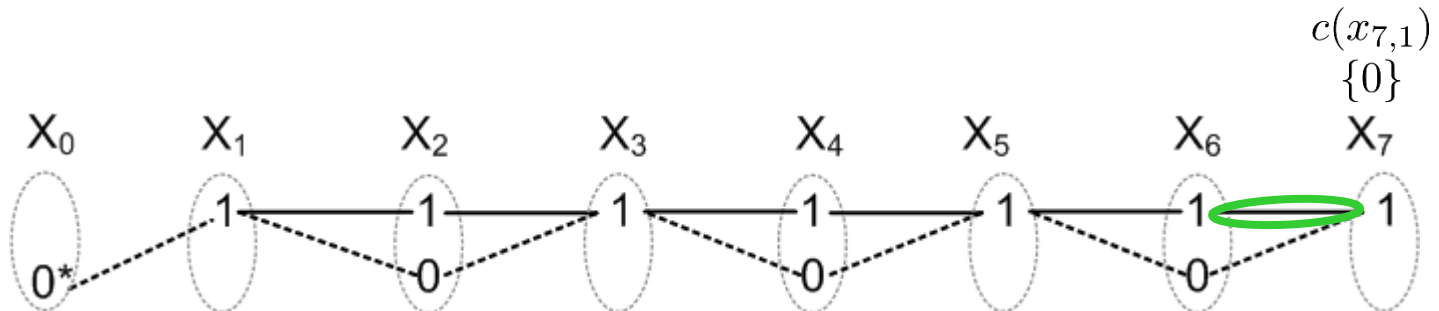
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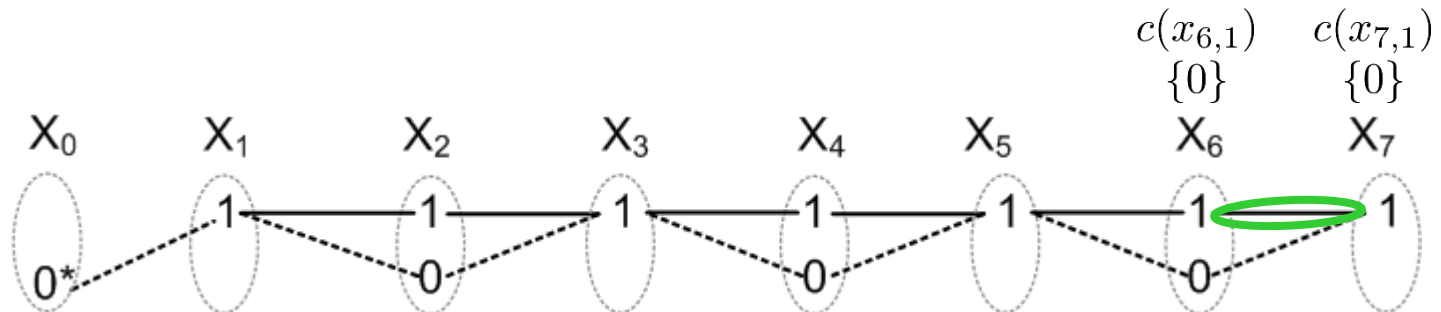
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SeqBin (DP-based propagator)

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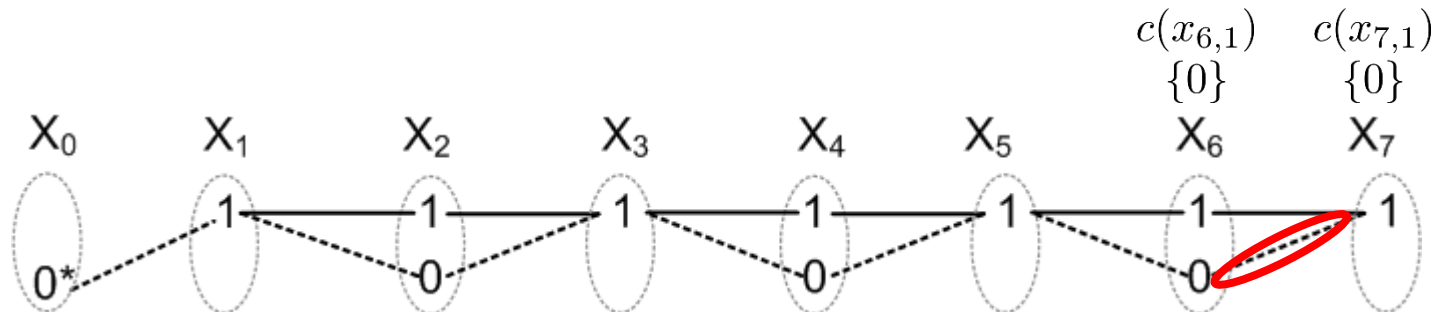
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SeqBin (DP-based propagator)

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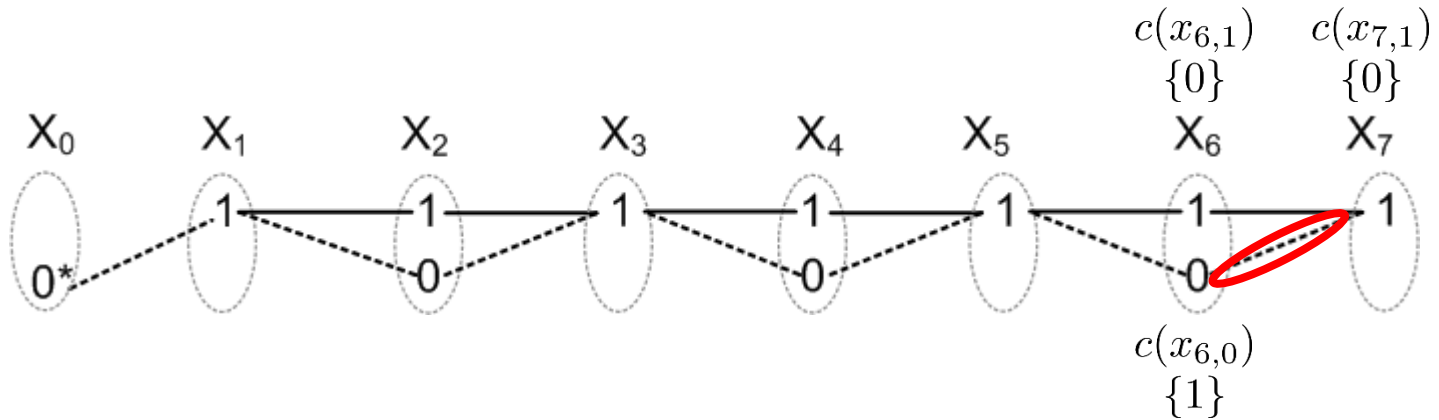
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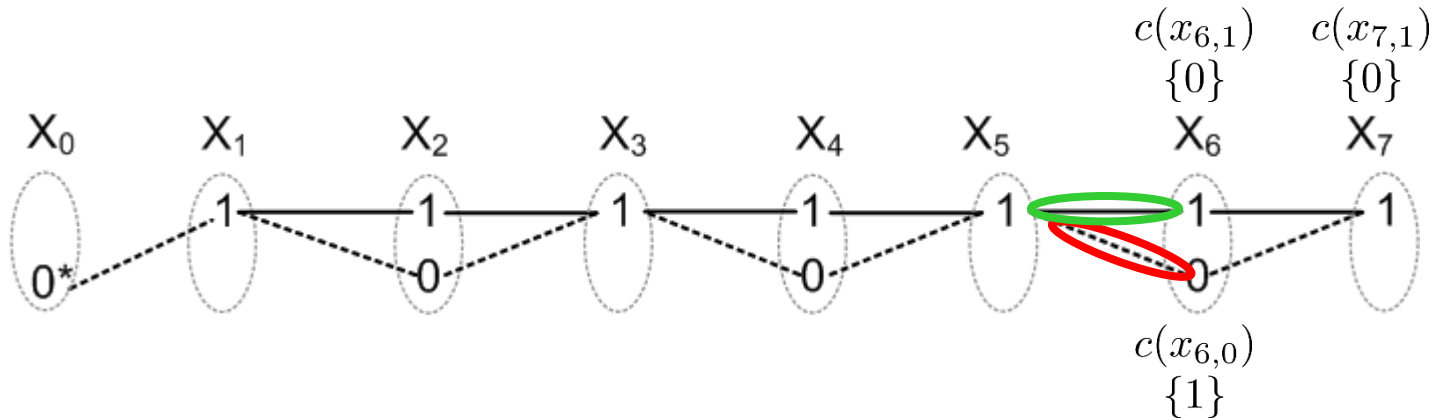
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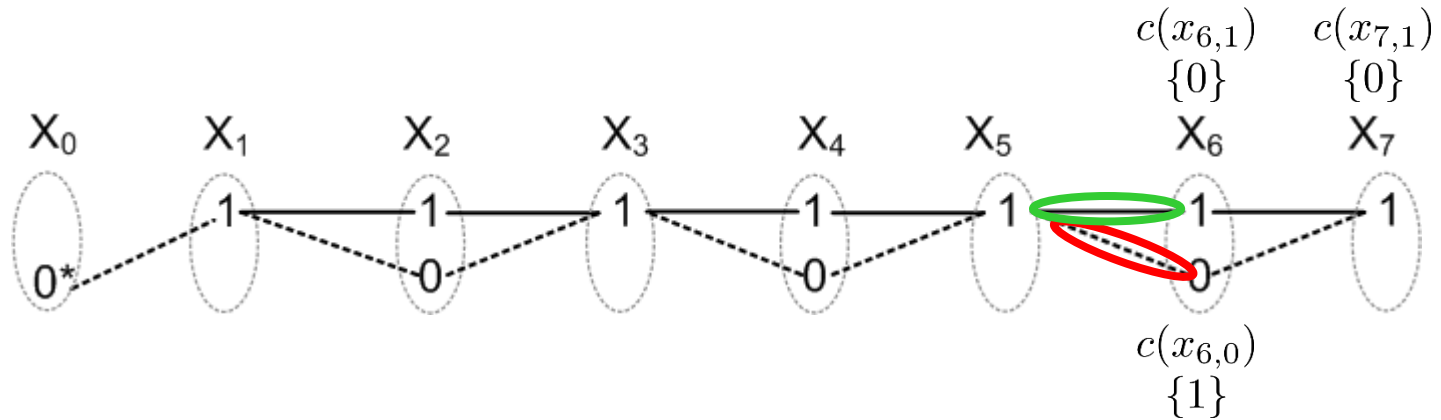
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SeqBin (DP-based propagator)

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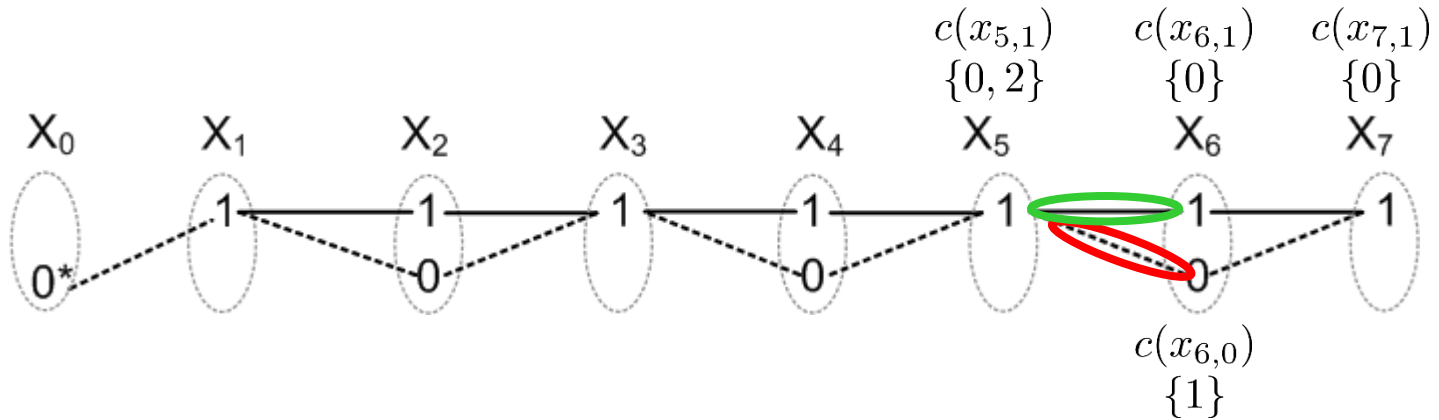


$$c(x_{5,1}) = (c(x_{6,1}) \uplus 0) \cup (c(x_{6,0}) \uplus 1)$$

SeqBin (DP-based propagator)

----- weight 1
 ——— weight 0

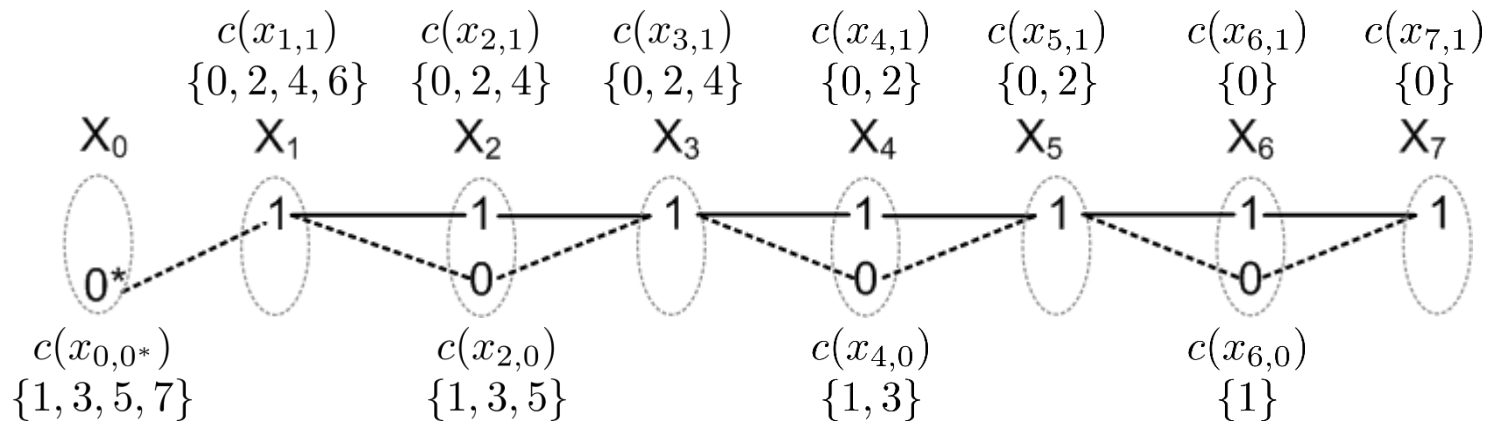
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SeqBin (DP-based propagator)

----- weight 1
 ——— weight 0

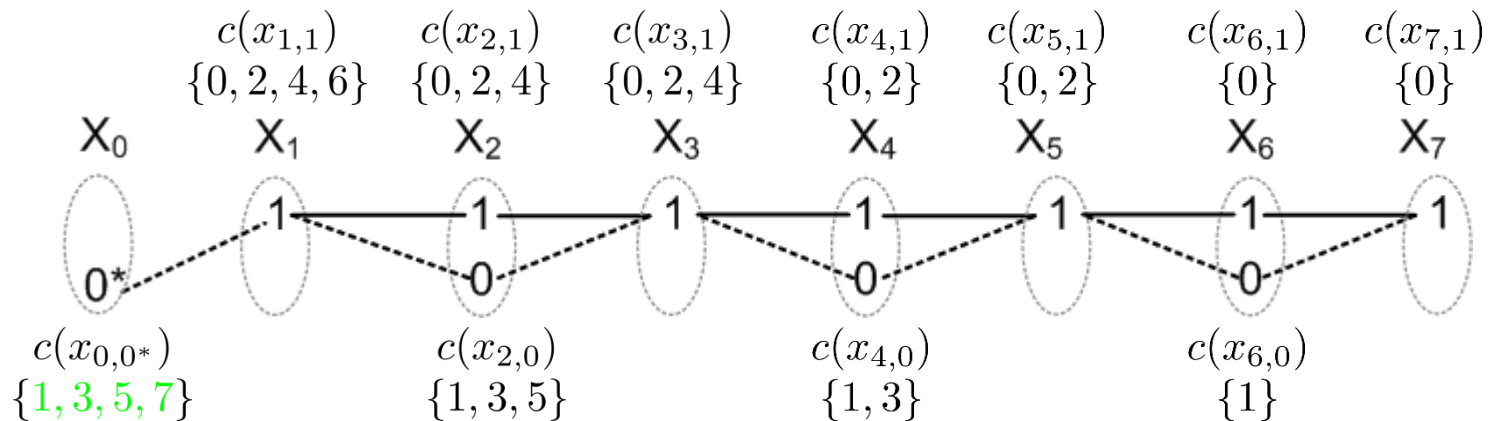
$$\text{SeqBin}(N, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$$



SeqBin (DP-based propagator)

----- weight 1
 ————— weight 0

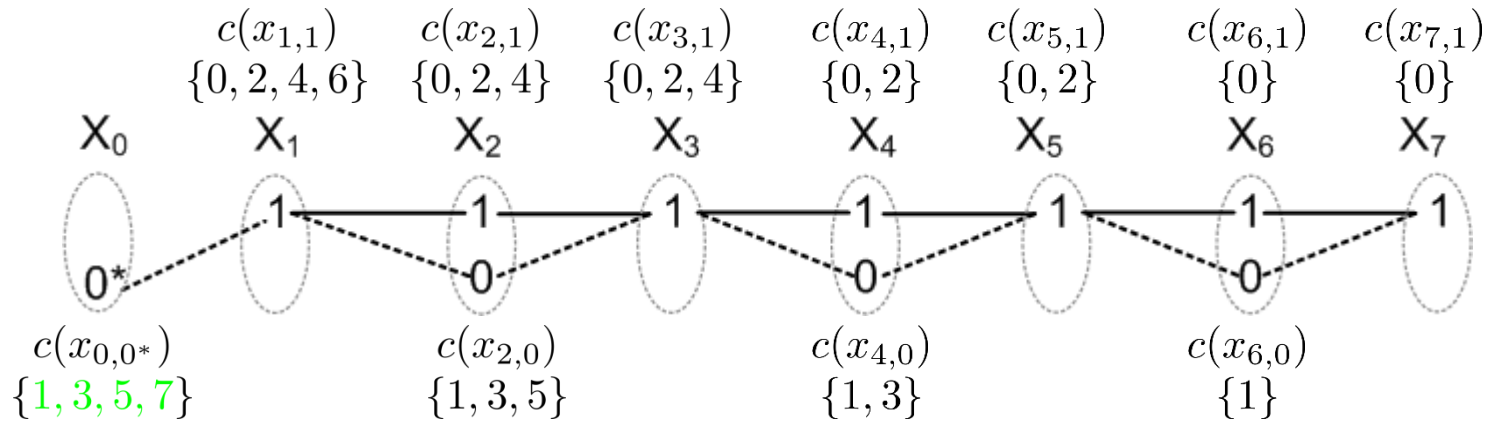
$$\text{SeqBin}(N, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True})$$



SeqBin (DP-based propagator)

----- weight 1
 ————— weight 0

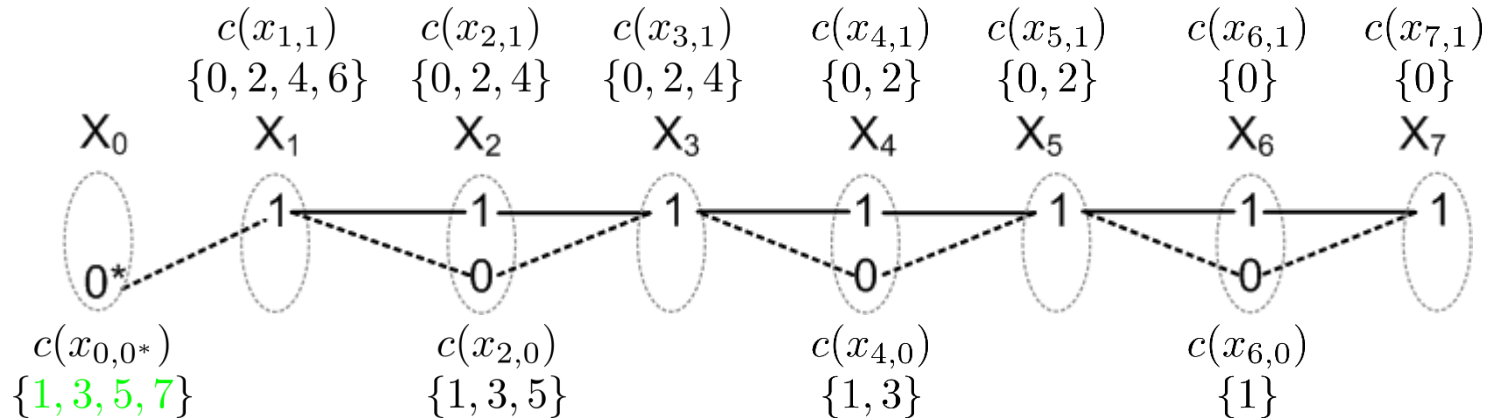
SeqBin($N = 3$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B = \text{True}$)



SeqBin (DP-based propagator)

----- weight 1
 ————— weight 0

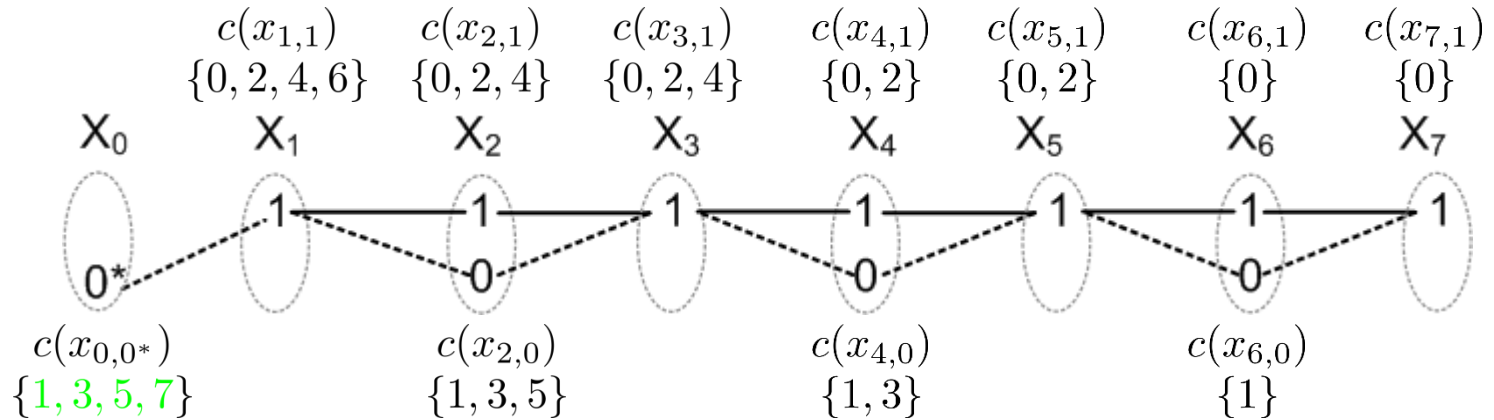
SeqBin($N = 4$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B = \text{True}$)



SeqBin (DP-based propagator)

----- weight 1
 ——— weight 0

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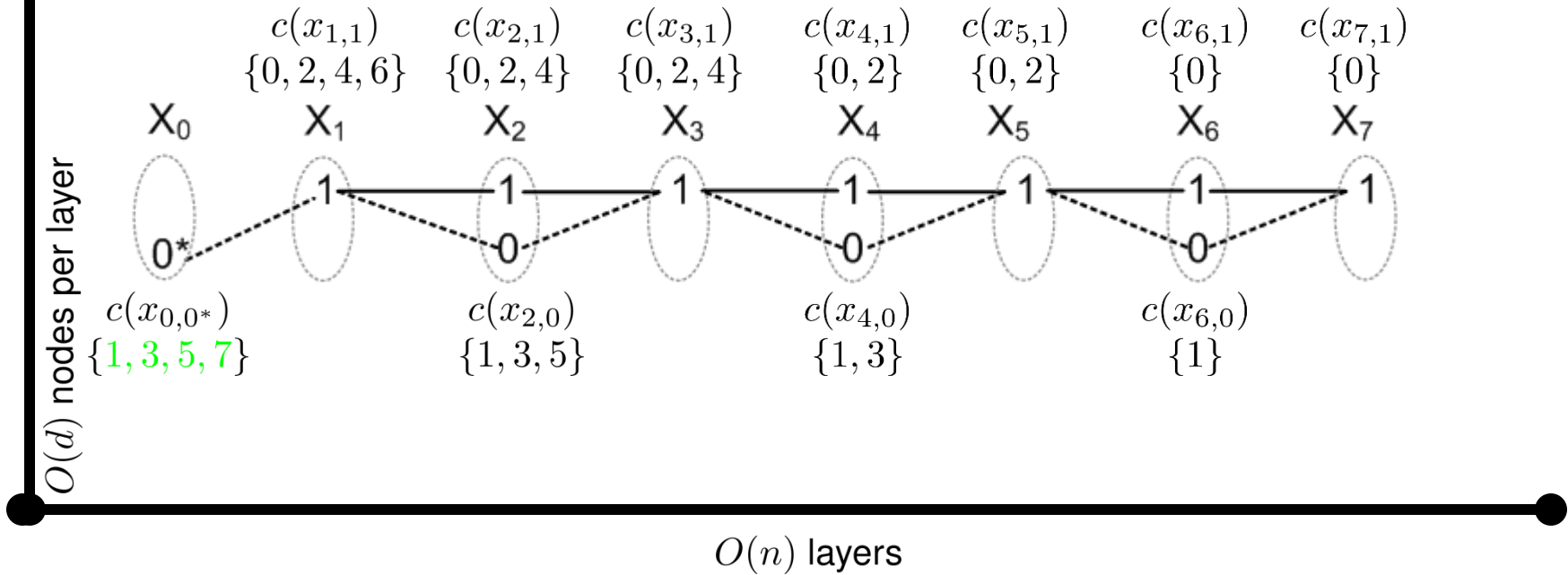


$O(n)$ layers

SeqBin (DP-based propagator)

----- weight 1
 ——— weight 0

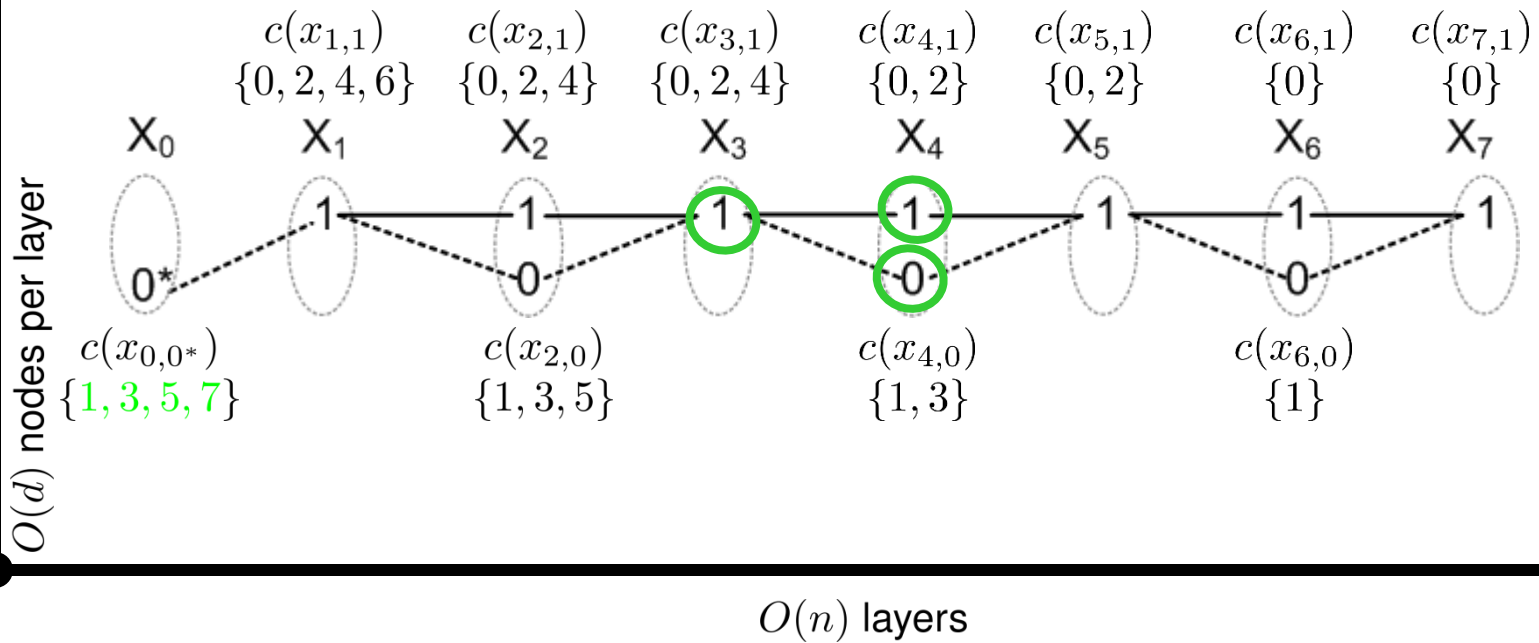
SeqBin($N = 4$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B = \text{True}$)



SeqBin (DP-based propagator)

----- weight 1
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SeqBin($N = 4$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B = \text{True}$)

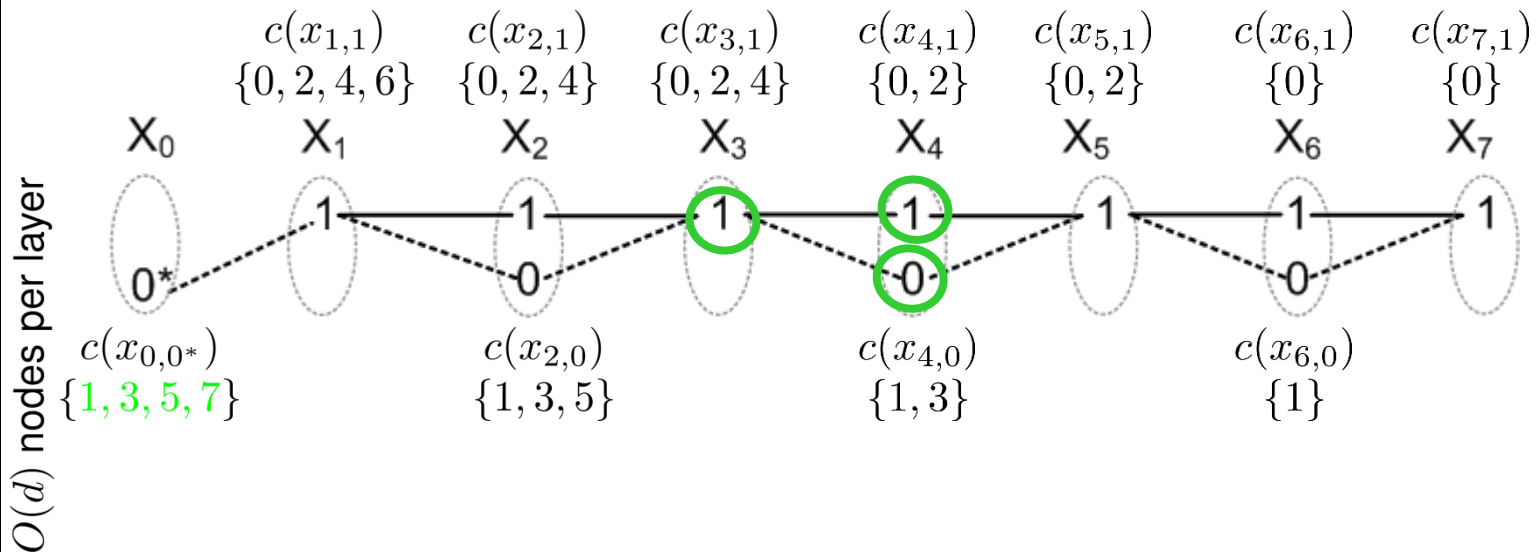


$$c(x_{3,1}) = (c(x_{4,1}) \uplus 0) \cup (c(x_{4,0}) \uplus 1) \quad O(nd)$$

SeqBin (DP-based propagator)

----- weight 1
 ——— weight 0

SeqBin($N = 4$, $[X_1, \dots, X_7]$, $C = \{(1, 1)\}$, $B = \text{True}$)



$$c(x_{i,j}) = \bigcup_{v \in D(X_{i+1})} (c(x_{i+1,v}) \uplus w(j, v)) \quad O(nd)$$

Related work

Related work
[IJCAI'11, Petit et al.]

Related work

$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$

Related work

the **monotone** constraint



$\text{SeqBin}(N, [X_1, \dots, X_7], C, B)$

Related work

Dynamic programming based filtering algorithm in $O(nd^2)$

Related work

$$c(x_{i,j}) = \bigcup_{v \in D(X_{i+1})} (c(x_{i+1,v}) \uplus w(j, v)) \quad O(n)$$

Related work

- does not detect bounds disentanglement
observed independently by Petit et al. , TR'2011

Related work

- does not detect bounds disentanglement
observed independently by Petit et al. , TR'2011
- not idempotent (i.e. $f(f(X)) \neq f(X)$)

Related work

Main problem: if B is monotone then

$$c(x_{i,j}) = [l, u]$$

Related work

Main problem: if B is monotone then

$$c(x_{i,j}) = [l, u]$$

$$c(x_{i,j}) = \bigcup_{v \in D(X_{i+1})} (c(x_{i+1,v}) \uplus w(j, v)) \quad O(d)$$

Related work

Main problem: if B is monotone then

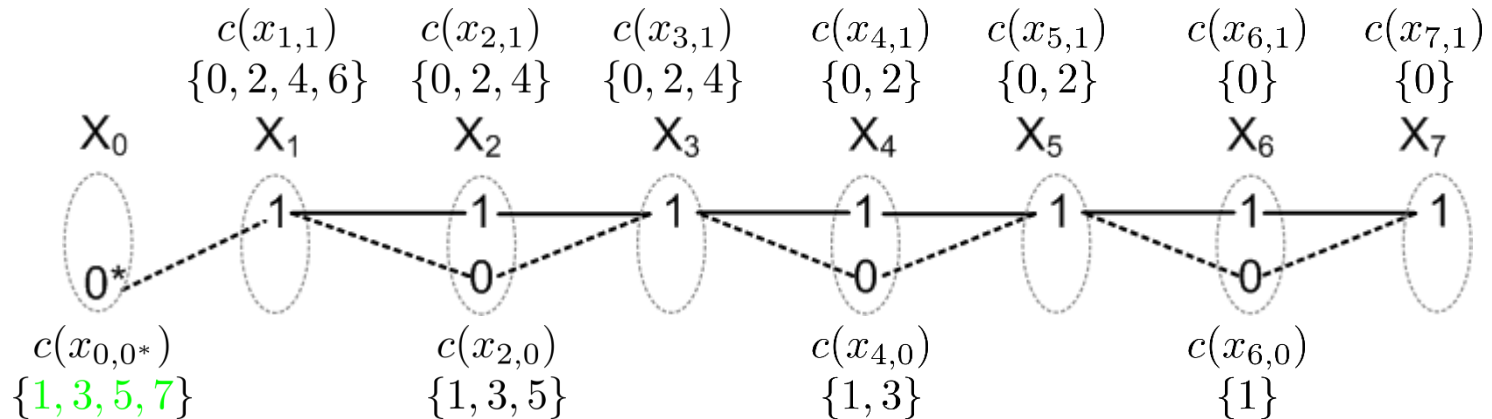
$$\cancel{c(x_{i,j}) = [l, u]}$$

$$c(x_{i,j}) = \bigcup_{v \in D(X_{i+1})} (c(x_{i+1,v}) \uplus w(j, v)) \quad O(d)$$

SeqBin (DP-based propagator)

----- weight 1
 ——— weight 0

SeqBin($N, [X_1, \dots, X_7], C = \{(1, 1)\}, B = \text{True}$)



Our contribution:
new DC filtering algorithm in $O(nd^2)$

New DC propagator

Goal: improve DP when B is monotone

New DC propagator



$$c(x_{i,j}) = \bigcup_{v \in D(X_{i+1})} (c(x_{i+1,v}) \uplus w(j, v)) \quad O(n)$$

Cost properties



$c(x_{i,j})$ is not an interval

Cost properties

 $c(x_{i,j})$ is not even an 'almost' interval 
(unbounded number of holes!)

Cost properties

$c(x_{i,j})$ has a special structure! 😄

Uniqueness

$c(x_{i,j})$ is a set of one of the following two forms:

zipper:

Uniqueness

$c(x_{i,j})$ is a set of one of the following two forms:

zipper: $\{2, 4, 6, 8\}$.

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i-zipper:

Uniqueness

$c(x_{i,j})$ is a set of one of the following two forms:

zipper: $\{2, 4, 6, 8\}$.

i-zipper: $[4, 11]$

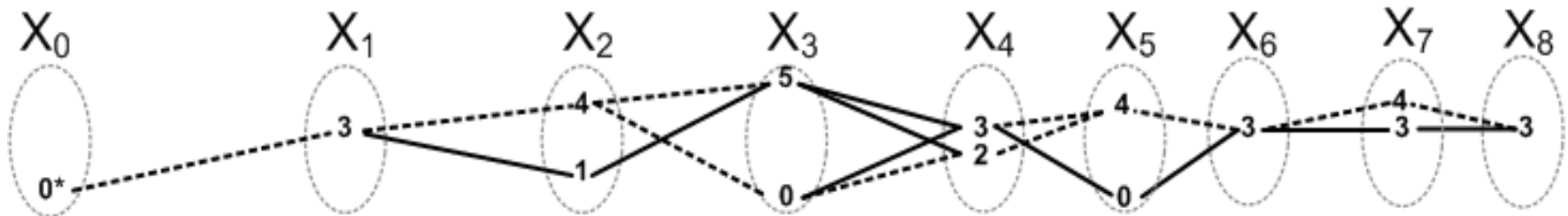
Uniqueness

$c(x_{i,j})$ is a set of one of the following two forms:

zipper: $\{2, 4, 6, 8\}$.

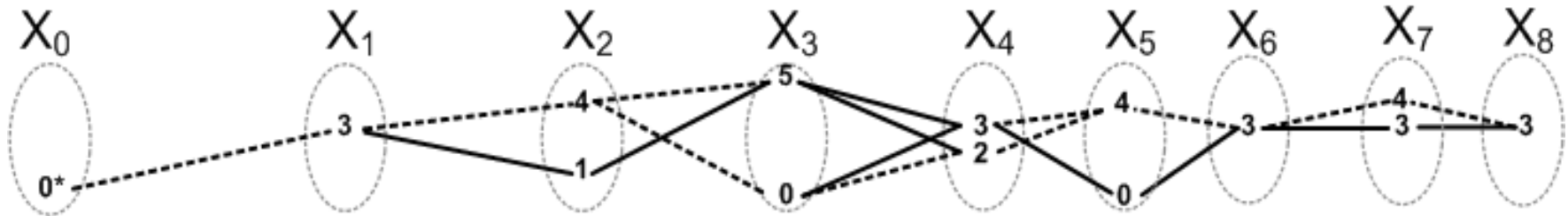
i-zipper: $\{2, 4\} \cup [4, 11] \cup \{11, 13, 15\}$

New DC propagator



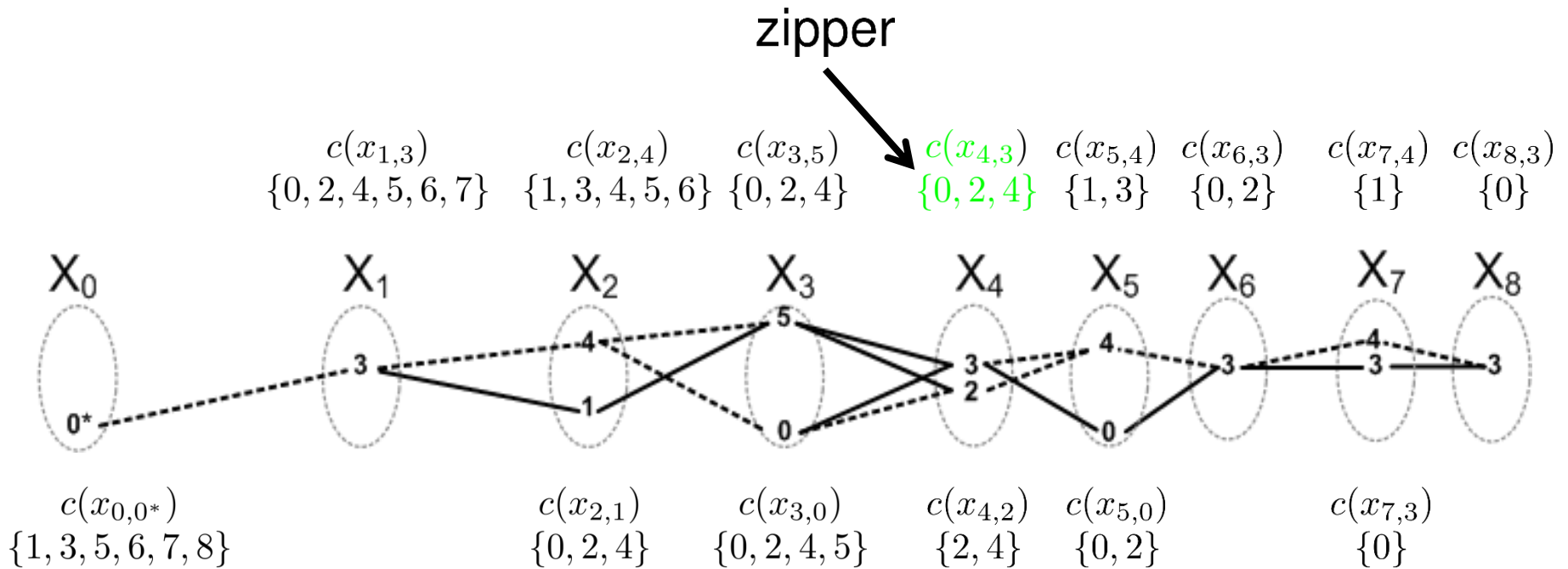
New DC propagator

	$c(x_{1,3})$	$c(x_{2,4})$	$c(x_{3,5})$	$c(x_{4,3})$	$c(x_{5,4})$	$c(x_{6,3})$	$c(x_{7,4})$	$c(x_{8,3})$
	{0, 2, 4, 5, 6, 7}	{1, 3, 4, 5, 6}	{0, 2, 4}	{0, 2, 4}	{1, 3}	{0, 2}	{1}	{0}

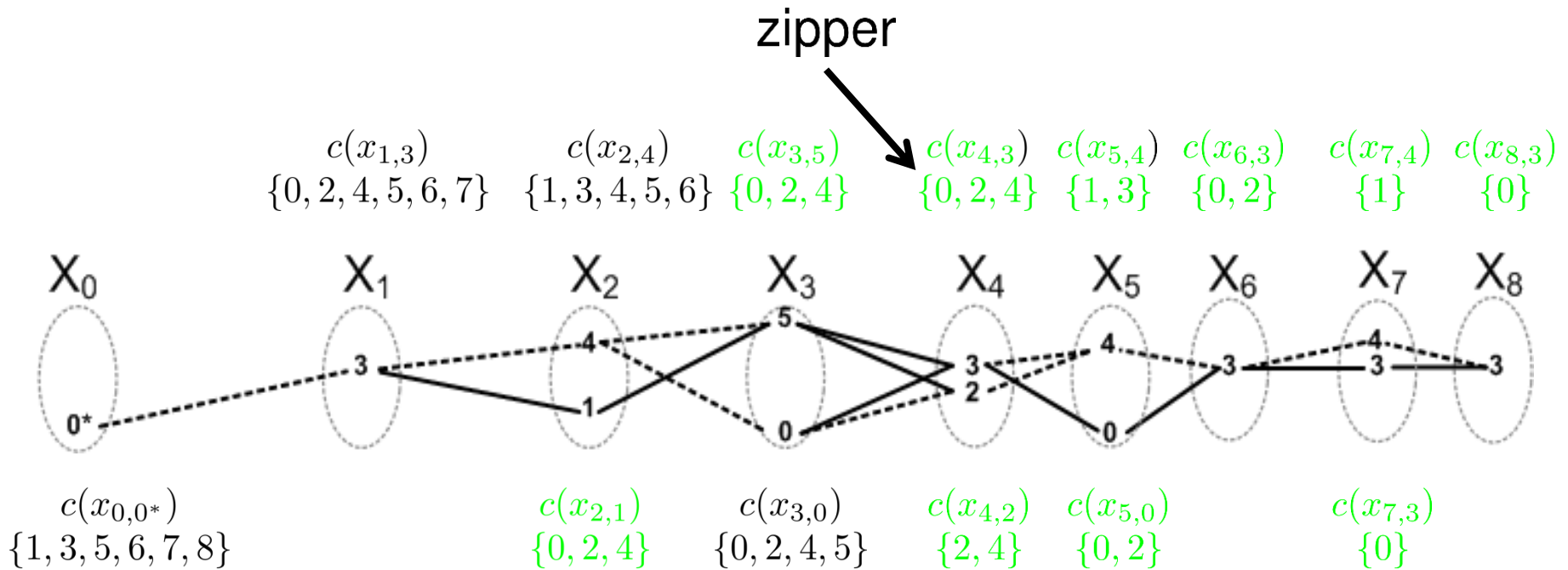


$c(x_{0,0^*})$	$c(x_{2,1})$	$c(x_{3,0})$	$c(x_{4,2})$	$c(x_{5,0})$	$c(x_{7,3})$
{1, 3, 5, 6, 7, 8}	{0, 2, 4}	{0, 2, 4, 5}	{2, 4}	{0, 2}	{0}

Uniqueness



Uniqueness

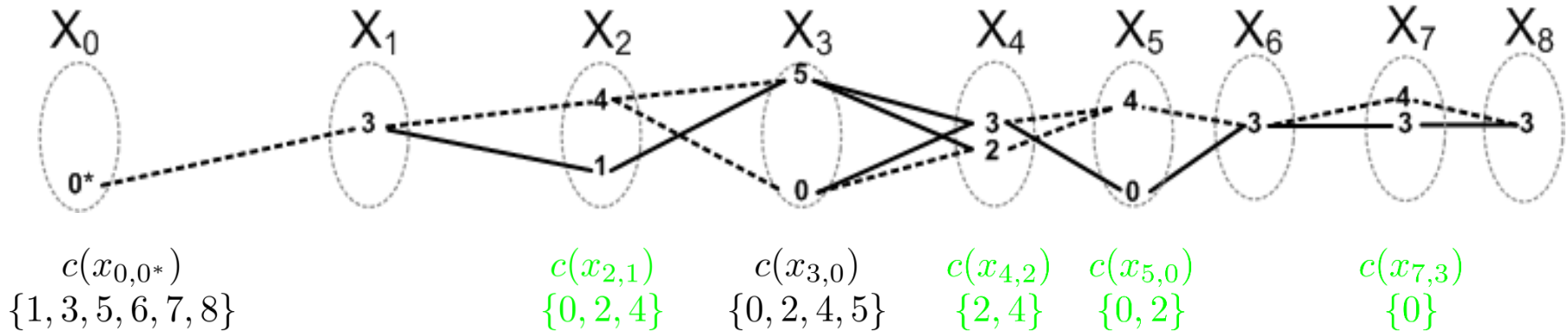


Uniqueness

i-zipper

zipper

$c(x_{1,3})$ $c(x_{2,4})$ $c(x_{3,5})$ $c(x_{4,3})$ $c(x_{5,4})$ $c(x_{6,3})$ $c(x_{7,4})$ $c(x_{8,3})$
 $\{0, 2, 4, 5, 6, 7\}$ $\{1, 3, 4, 5, 6\}$ $\{0, 2, 4\}$ $\{0, 2, 4\}$ $\{1, 3\}$ $\{0, 2\}$ $\{1\}$ $\{0\}$

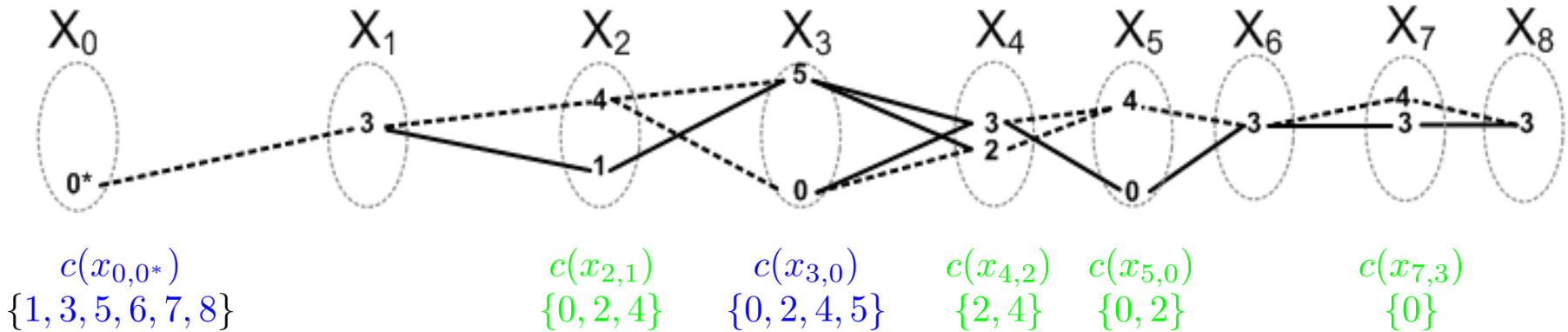


Uniqueness

i-zipper

zipper

$c(x_{1,3})$ $c(x_{2,4})$ $c(x_{3,5})$ $c(x_{4,3})$ $c(x_{5,4})$ $c(x_{6,3})$ $c(x_{7,4})$ $c(x_{8,3})$
 $\{0, 2, 4, 5, 6, 7\}$ $\{1, 3, 4, 5, 6\}$ $\{0, 2, 4\}$ $\{0, 2, 4\}$ $\{1, 3\}$ $\{0, 2\}$ $\{1\}$ $\{0\}$



Structure (Bounded holes)

The number of holes in an i-zipper is bounded.

Structure (Bounded holes)

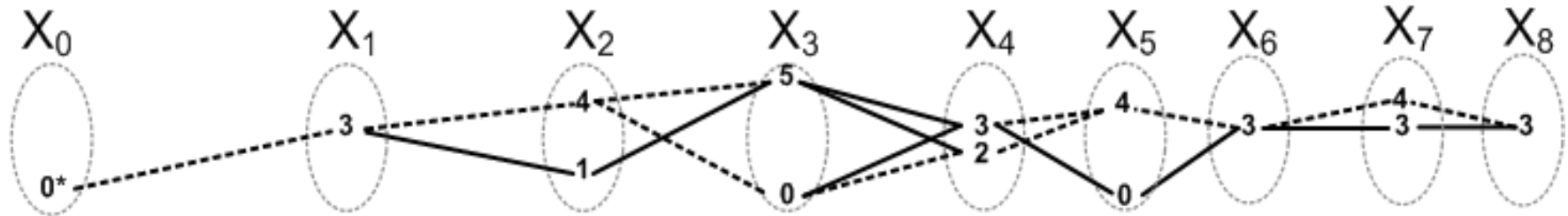
i-zipper: $\{2, 4\} \cup [4, 11] \cup \{11, 13, 15\}$

Structure (Closeness)

Bounds of $c(x_{i,j})$ and $c(x_{i,k})$ are close to each other.

Structure (Closeness)

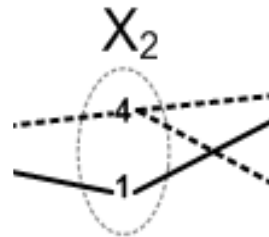
	$c(x_{1,3})$	$c(x_{2,4})$	$c(x_{3,5})$	$c(x_{4,3})$	$c(x_{5,4})$	$c(x_{6,3})$	$c(x_{7,4})$	$c(x_{8,3})$
	{0, 2, 4, 5, 6, 7}	{1, 3, 4, 5, 6}	{0, 2, 4}	{0, 2, 4}	{1, 3}	{0, 2}	{1}	{0}



$c(x_{0,0^*})$	$c(x_{2,1})$	$c(x_{3,0})$	$c(x_{4,2})$	$c(x_{5,0})$	$c(x_{7,3})$
{1, 3, 5, 6, 7, 8}	{0, 2, 4}	{0, 2, 4, 5}	{2, 4}	{0, 2}	{0}

Structure (Closeness)

$$c(x_{2,4})$$
$$\{1, 3, 4, 5, 6\}$$



$$c(x_{2,1})$$
$$\{0, 2, 4\}$$

More properties

and many more properties...

How did we know?

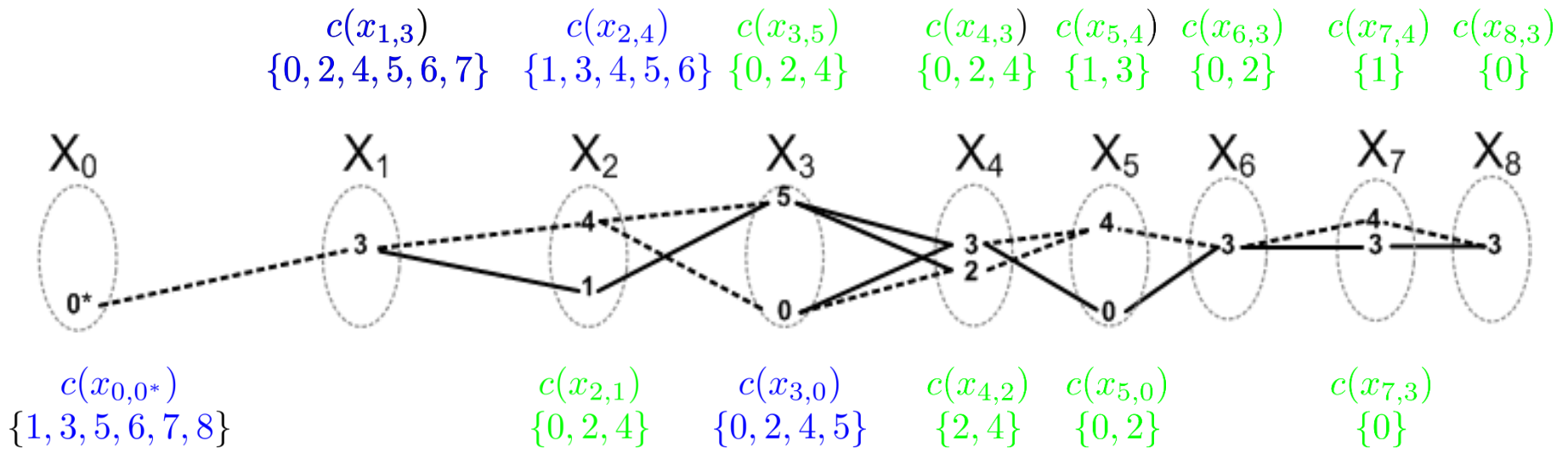
zipper: $\{2, 4, 6, 8\}$.

i-zipper: $\{2, 4\} \cup [4, 11] \cup \{11, 13, 15\}$

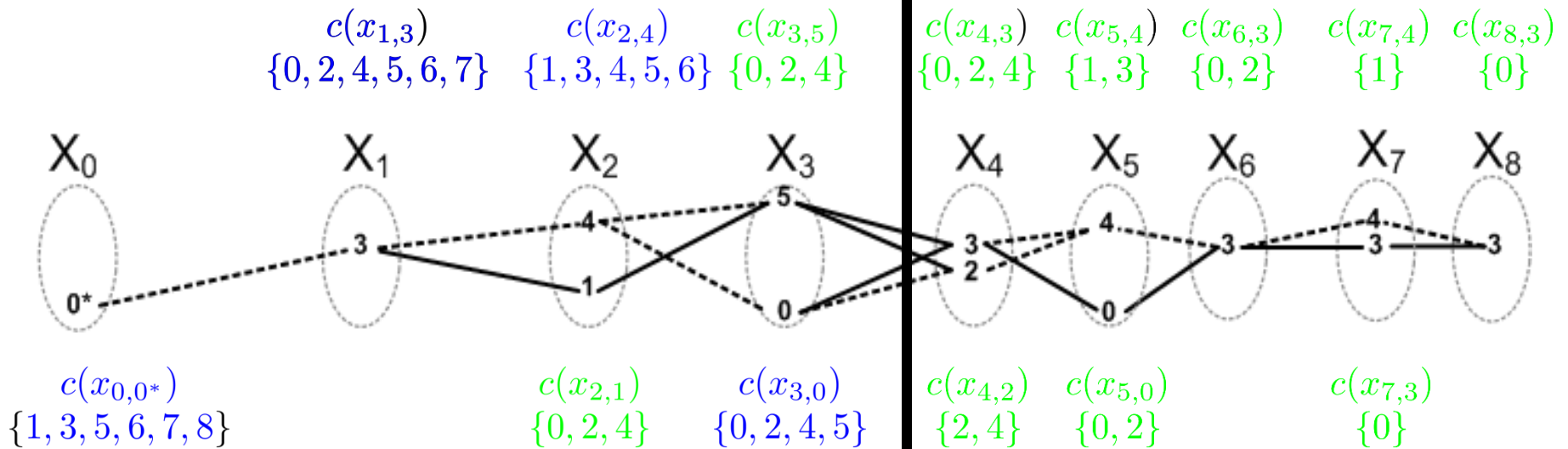
SeqBin

Global structure of cost

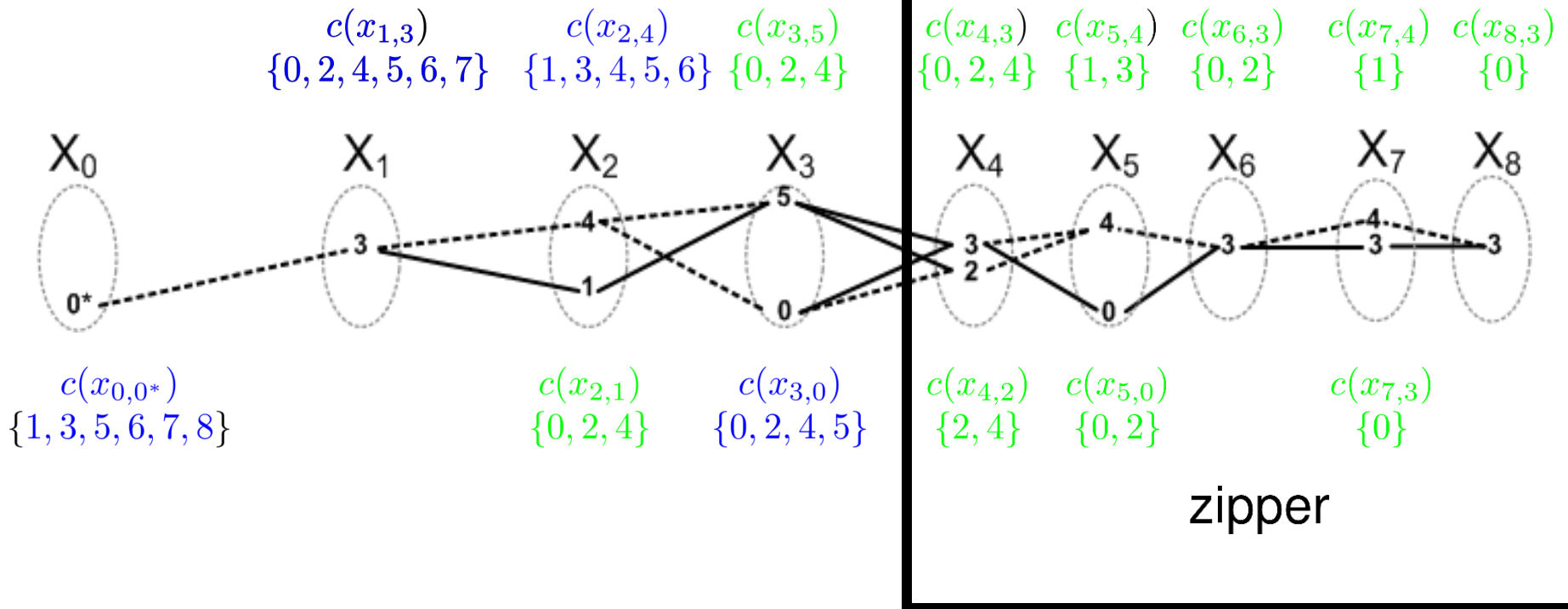
Blocks



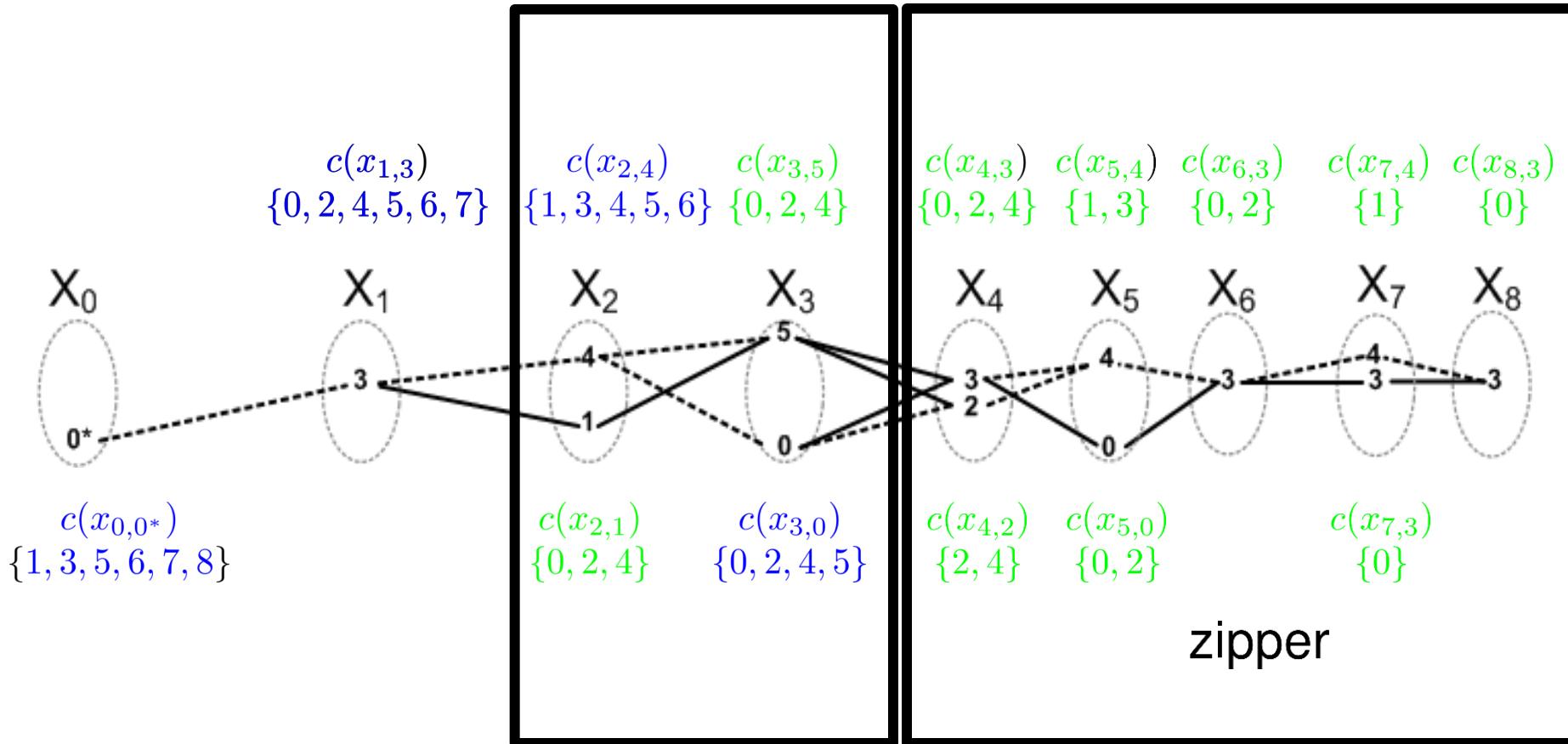
Blocks



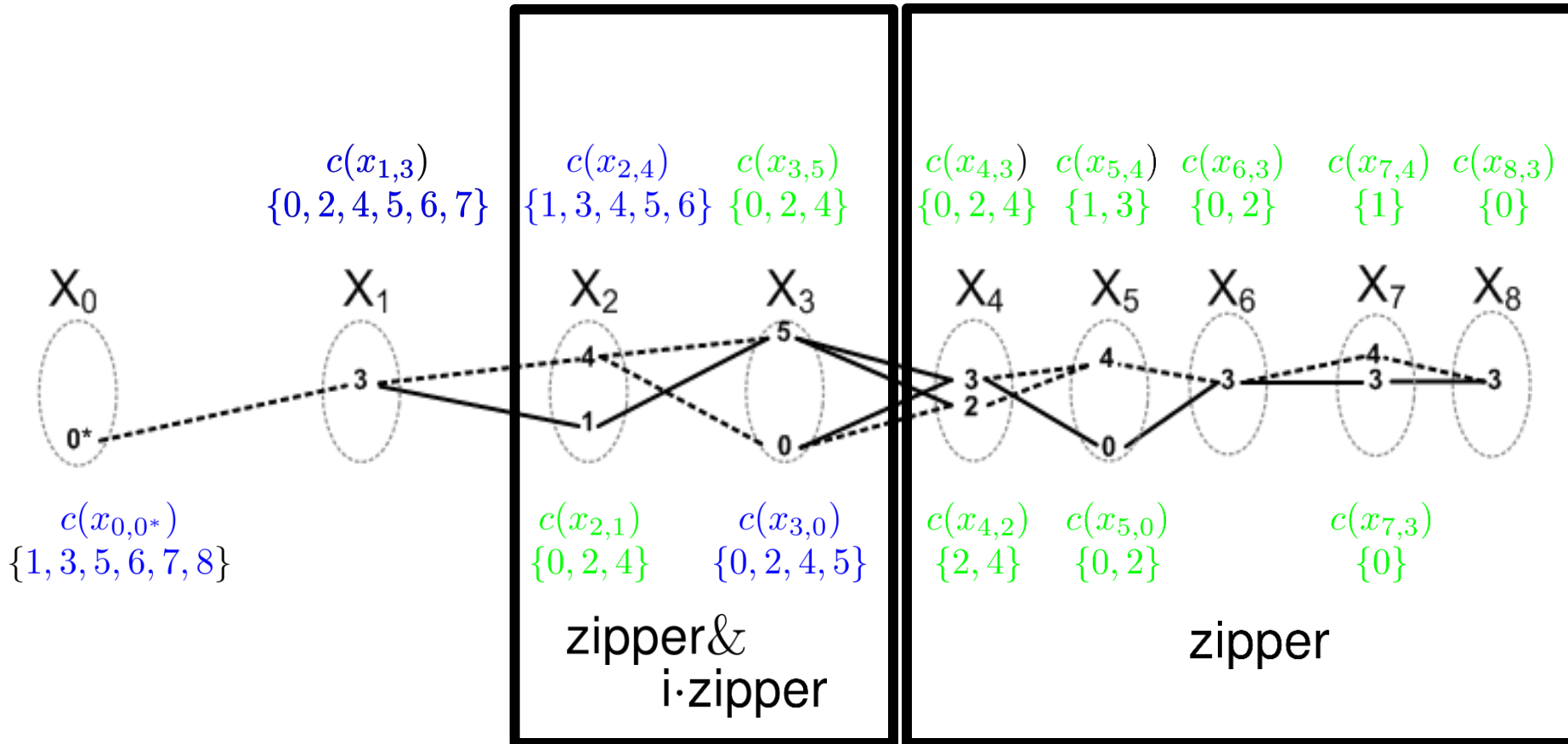
Blocks



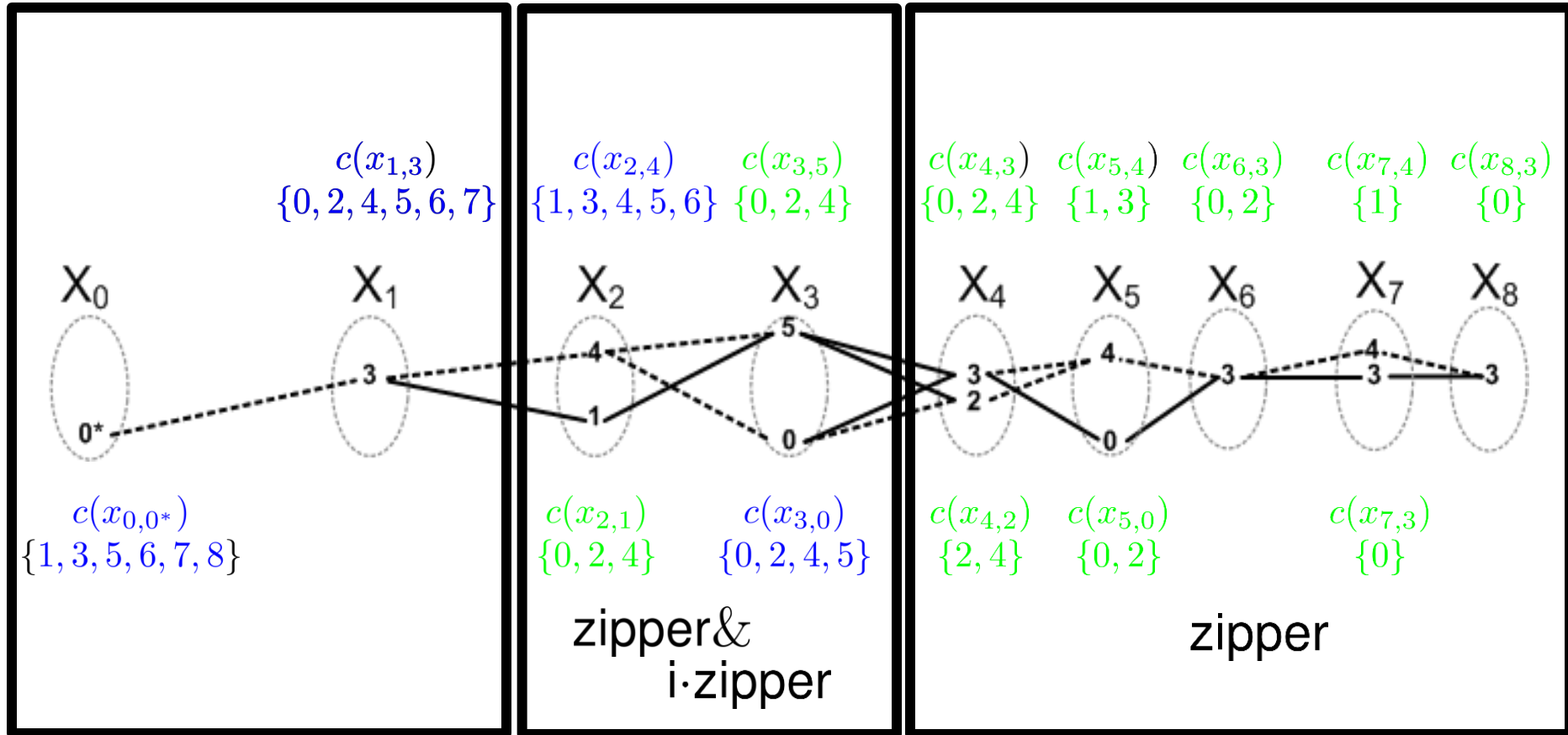
Blocks



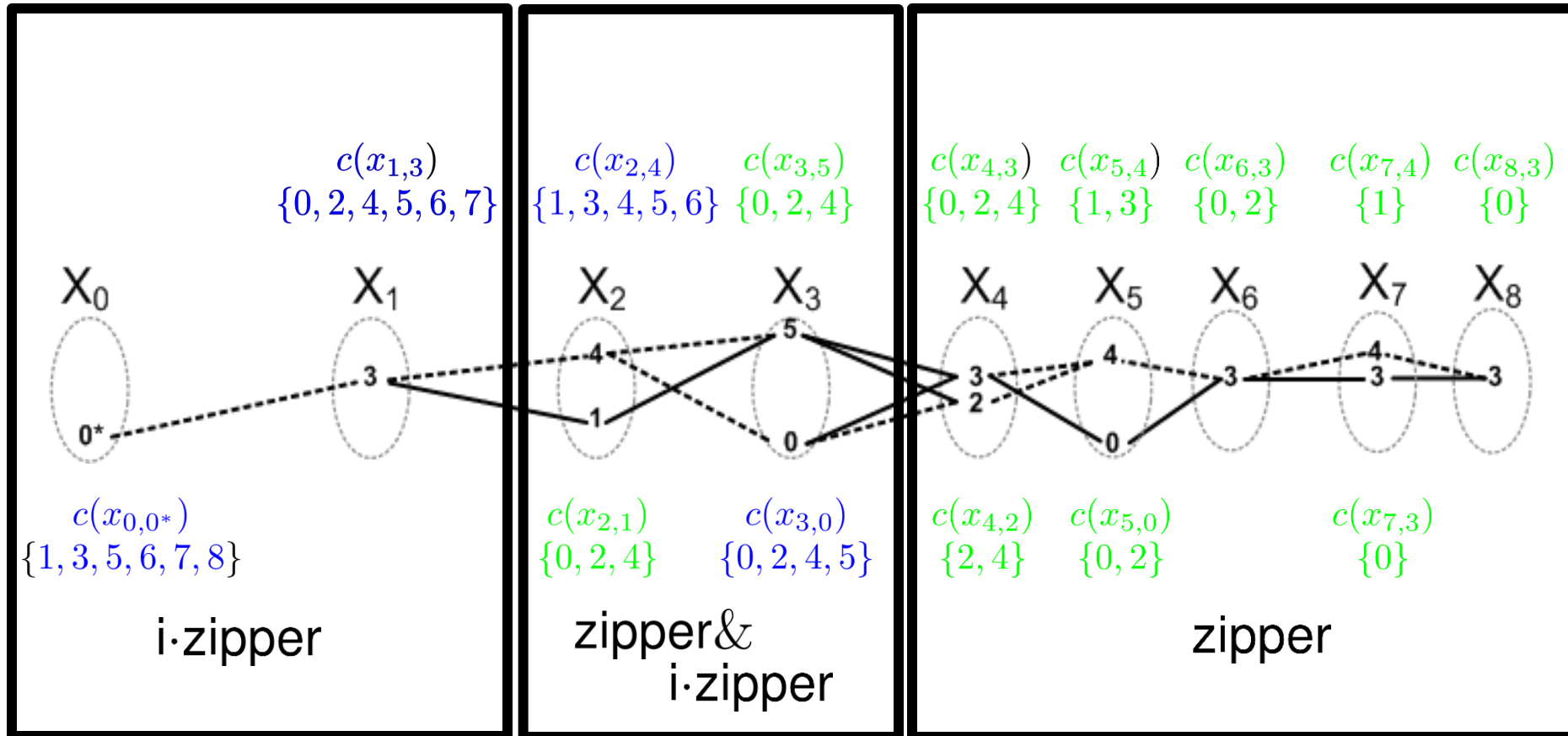
Blocks



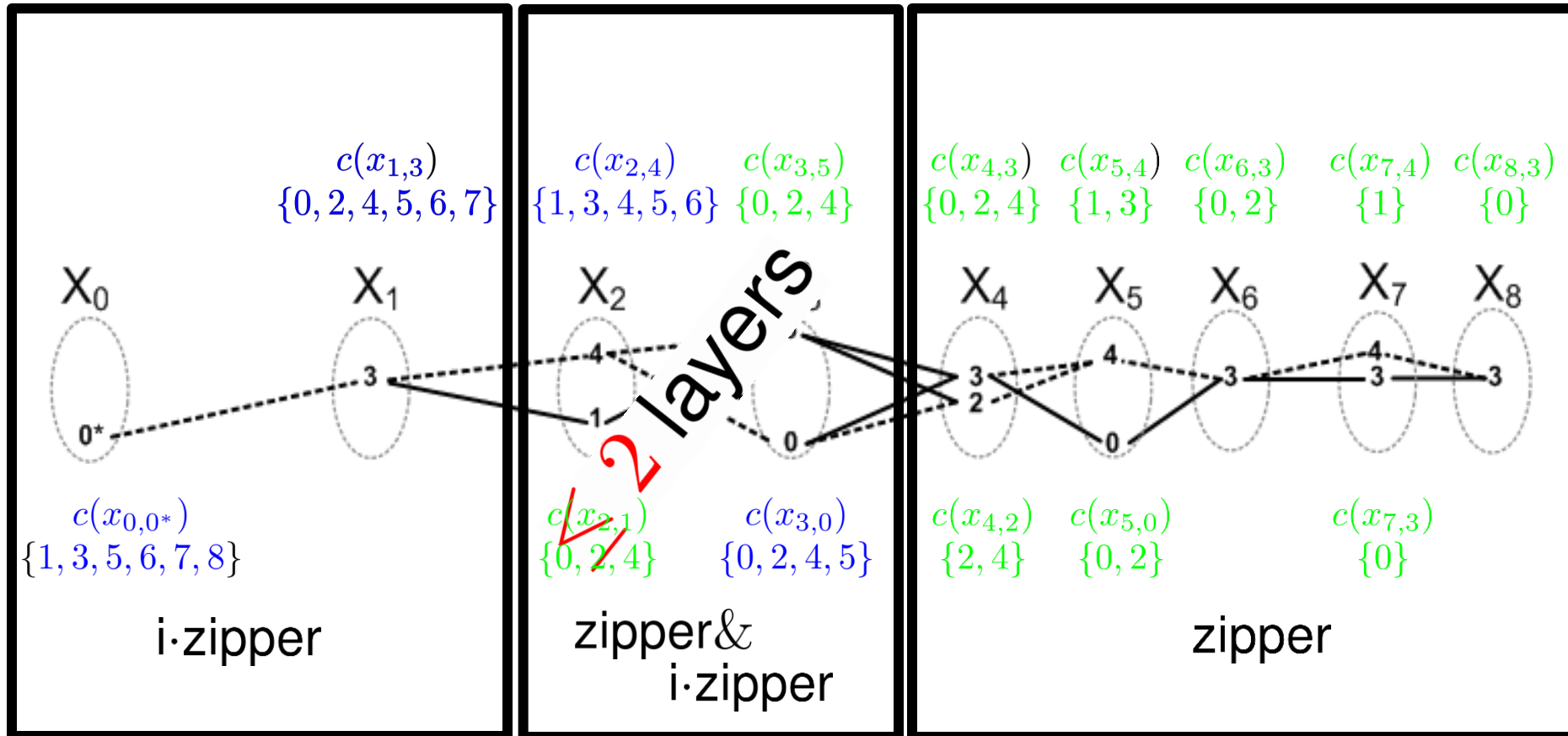
Blocks



Blocks



Blocks



Conclusions

Contributions

If B is monotone then DC propagator runs in $O(nd^2)$

Contributions

If B is monotone then DC propagator runs in $O(nd^2)$ 😄

If B is monotone and C is convex

then DC propagator runs in $O(nd)$ 😄

Thank you!